

A two-symbol system of encoding and some of its applications

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ABSTRACT. The article is devoted to a two-symbol system of encoding for real numbers with two bases of different signs $g_0 \in (0; \frac{1}{2}]$ and $g_1 \equiv g_0 - 1$, as well as its applications in metric number theory and the metric theory of functions. We prove that any natural number a can be represented as

$$a = 2^n + \sum_{k=1}^n [(-1)^{1+\sigma_k} a_k 2^{n-k}] \equiv (1a_1 \dots a_n)_G,$$

where $a_k \in \{0; 1\}$ and $\sigma_k = a_1 + \dots + a_{k-1}$, and there exist exactly two such representations. Any number $x \in (0; g_0]$ can be represented as

$$\delta_{\alpha_1} + \sum_{k=2}^{\infty} (\delta_{\alpha_k} \prod_{j=1}^{k-1} g_{\alpha_j}) \equiv \Delta_{\alpha_1 \alpha_2 \dots \alpha_k \dots}^{G_2}, \delta_{\alpha_k} = \alpha_k g_{1-\alpha_k}.$$

Most numbers have a unique G_2 -representation, while a countable set has exactly two representations: $\Delta_{c_1 \dots c_m 01(0)}^{G_2} = \Delta_{c_1 \dots c_m 11(0)}^{G_2}$.

For $g_0 = \frac{1}{2}$, any number x in the interval $[0; 1]$ has the expansion

$$x = \frac{1}{2} \alpha_0 + \sum_{k=1}^{\infty} \frac{\alpha_k (-1)^{1+\sigma_k}}{2^k} \equiv \Delta_{\alpha_0 \alpha_1 \dots \alpha_n \dots}.$$

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Introduction

Today, exact sciences confidently operate with both finite and infinite sets and data arrays. The ideas of coordinate representation and encoding are effectively applied. As a result, powerful families of mathematical objects are described using a small number of basis (reference) objects and relations. Meaningful information about dependencies and correspondences takes a digital form, encoded by codes: sets, matrices, and sequences of elements of alphabet, which can be finite or infinite, constant or variable [11].

Traditionally, two-symbol systems of encoding of information use the alphabet $A = \{0, 1\}$. So far, they remain unmatched in applications, although three-symbol systems are more efficient in some sense. Two-symbol systems deserve a special attention, particularly due to the minimality of their alphabet. In the sequel, we will focus on the encoding of real numbers.

An encoding of real numbers from the set D using the alphabet A is defined as a surjective mapping (onto mapping) from the space $L = A \times A \times \dots$ of sequences of elements of the alphabet A onto the set D . Traditionally, the binary numeral system is the simplest two-symbol encoding of numbers. Its metric counterpart is the negabinary numeral system [6], a system with a negative base (-2) .

This work is devoted to an encoding system that is fundamentally different from the above-mentioned systems.

Such representation of numbers in the interval $[0; g_0]$ and its various applications were studied in the papers [4, 6–10].

Except for [1–3, 5], we are not aware of any other papers where numeral systems with two bases are considered.

1. Existence of a G -representation of a natural number

Definition 1. If, for a natural number a , there exists a sequence $(a_1, \dots, a_n)$ of zeros and ones such that

$$a = 2^n + \sum_{k=1}^n [(-1)^{1+\sigma_k} a_k 2^{n-k}],$$

where $\sigma_k = a_1 + \dots + a_{k-1}$, then we say that the number a has a G -representation. Symbolically, this is written as

$$a = (1a_1 \dots a_n)_G.$$

The number a is thus an $(n + 1)$ -digit number.

For example, each number within the range from 1 to 10 has a G -representation, and it has exactly two representations. It is easy to verify this statement:

$$\begin{aligned}
 1 &= 2^0 = (1)_G = 2^1 - 2^0 = (11)_G, \\
 2 &= 2^1 = (10)_G = 2^2 - 2^1 = (110)_G, \\
 3 &= 2^2 - 2^1 + 2^0 = (111)_G = 2^2 - 2^0 = (101)_G, \\
 4 &= 2^2 = (100)_G = 2^3 - 2^2 = (1100)_G, \\
 5 &= 2^3 - 2^2 + 1 = (1101)_G = 2^3 - 2^2 + 2 - 1 = (1111)_G, \\
 6 &= 2^3 - 2^2 + 2 = (1110)_G = 2^3 - 2 = (1010)_G, \\
 7 &= 2^3 - 1 = (1001)_G = 2^3 - 2 + 1 = (1011)_G, \\
 8 &= 2^3 = (1000)_G = 2^4 - 2^3 = (11000)_G, \\
 9 &= 2^4 - 2^3 + 1 = (11001)_G = 2^4 - 2^3 + 2 - 1 = (11011)_G, \\
 10 &= 2^4 - 2^3 + 2 = (11010)_G = 2^4 - 2^3 + 2 - 1 = (11110)_G.
 \end{aligned}$$

Remark 1. As we can see, different G -representations of the same natural number can have a different number of digits. For example, $8 = (1000)_G = (11000)_G$, and in general

$$a = 2^n = (1 \underbrace{0 \dots 0}_n)_G = (11 \underbrace{0 \dots 0}_n)_G,$$

$$\text{while } a = 2^n + 1 = 2^{n+1} - 2^n + 1 = (11 \underbrace{0 \dots 01}_{n-1})_G = (11 \underbrace{0 \dots 011}_{n-2})_G.$$

Theorem 1. *Each natural number has exactly two G -representations, that is, for any natural number a , there exists a sequence of zeros and ones $(a_1, a_2, \dots, a_n)$ such that*

$$a = 2^n + \sum_{k=1}^n [(-1)^{1+\sigma_k} a_k 2^{n-k}] \equiv (1a_1 \dots a_n)_G, \quad (1)$$

where $\sigma_1 = 0$, $\sigma_k = a_1 + \dots + a_{k-1}$, and there are exactly two such expansions.

Proof. Obviously, for any natural number a , there exists a natural number n such that $2^{n-1} < a \leq 2^n$. We will use the method of mathematical induction on n . For $n = 1$, we have $1 < a \leq 2$ and the statement is evident for the numbers 1 and 2 (see the example above).

Assume that the statement holds for $n = k$, i.e., for $a \in (2^{k-1}; 2^k]$, there exists an expansion $a = 2^k + \sum_{i=1}^k [(-1)^{1+\sigma_i} a_i 2^{k-i}] \equiv (1a_1 \dots a_k)_G$, where $\sigma_i = a_1 + \dots + a_{i-1}$, and there are exactly two such expansions.

Consider $n = k + 1$, i.e., the case of a number $a \in (2^k; 2^{k+1}]$. If $a = 2^{k+1}$, then $a = (\underbrace{10 \dots 0}_{k+1})_G = (\underbrace{110 \dots 0}_{k+1})_G$. Now, suppose that $2^k < a < 2^{k+1}$; then $0 < a - 2^k \equiv u < 2^k$. By the induction hypothesis, the number u has a G -representation $u = (1c_1 c_2 \dots c_m)_G$, where $m < k$. Thus, we have $a = 2^k + u = 2^{k+1} - 2^k + u = (\underbrace{110 \dots 0}_{k-m} c_1 \dots c_m)_G$. Since the number $u \equiv a - 2^k$ satisfies the inequality $0 < u < 2^k$, it has exactly two G -representations by the induction hypothesis. Therefore, the number a also has exactly two G -representations. By the principle of mathematical induction, the statement is proven for any natural number a . $\square$

Definition 2. We say that a G -representation of a natural number is called *canonical* if it has the minimum number of digits and the maximum number of zeros simultaneously.

For example, the canonical G -representation of the numbers

$$16 = (10000)_G = (110000)_G, \quad a = 2^n - 1 = (\underbrace{10 \dots 01}_{n-1})_G = (\underbrace{10 \dots 011}_{n-2})_G$$

is the first of the G -representations.

2. Identification and comparison of natural numbers

Theorem 2. If $a = (1a_1 \dots a_n a_{n+1})_G$ and $b = (1b_1 \dots b_n)_G$, then $a \geq b$, where equality holds only when $a = (\underbrace{110 \dots 0}_m)_G$ and $b = (\underbrace{10 \dots 0}_n)_G$. If the number of digits in the G -representation of $c = (1c_1 \dots c_n \dots c_{n+p})_G$ exceeds the number of digits in the G -representation of b by p , then $c > b$.

Proof. We only need to consider the case where one number has exactly one more digit than the other. Let $a = (1a_1 \dots a_n a_{n+1})_G$ and $b = (1b_1 \dots b_n)_G$.

Consider the difference

$$\rho \equiv a - b = (2^{n+1} - 2^n) + \left(\sum_{i=1}^{n+1} [(-1)^{1+\sigma_i} a_i 2^{n+1-i}] - \sum_{i=1}^n [(-1)^{1+\sigma'_i} b_i 2^{n-i}] \right).$$

Since $\min \sum_{i=1}^{n+1} [(-1)^{1+\sigma_i} a_i 2^{n+1-i}] = -2^n$ is attained only when $1 - a_1 = a_2 = a_3 = \dots = a_{n+1} = 0$, and $\max \sum_{i=1}^n [(-1)^{1+\sigma'_i} b_i 2^{n-i}] = 0$ is attained only when $b_1 = b_2 = \dots = b_n = 0$, it follows that $\rho \geq 2^{n+1} - 2^n - 2^n = 0$, where equality holds only when $a = (11\underbrace{0\dots 0}_m)_G$ and $b = (1\underbrace{0\dots 0}_n)_G$. By the theorem assumption, let $c = (1c_1\dots c_n\dots c_{n+p})_G$, where $p \geq 2$, then the difference

$$\begin{aligned} \delta &= c - b = (2^{n+p} - 2^n) + \\ &+ \left(\sum_{i=1}^{n+p} [(-1)^{1+\sigma_i} c_i 2^{n+1-i}] - \sum_{i=1}^n [(-1)^{1+\sigma'_i} b_i 2^{n-i}] \right) \geq \\ &\geq 2^{n+p} - 2^n - 2^{n+p-1} + 0 = 2^n(2^{p-1} - 1) > 0, \end{aligned}$$

which implies $c > b$. This completes the proof. $\square$

Remark 2. Two G -representations of a natural number $a \neq 2^n$ have the same number of digits, while the G -representations of numbers of the form 2^n have different numbers of digits.

Theorem 3. *The numbers*

$$a = (1a_1\dots a_{k-1}1a_{k+1}\dots a_n)_G \quad \text{and} \quad b = (1a_1\dots a_{k-1}0b_{k+1}\dots b_n)_G$$

satisfy the following relation:

- 1) $a \geq b$ if σ_k is odd;
- 2) $a \leq b$ if σ_k is even, where equality holds only if

$$1 - a_{k+1} = a_{k+2} = \dots = a_n = 0 = 1 - b_{k+1} = b_{k+2} = \dots = b_n.$$

Proof. Since $a_k = 1$ and $b_k = 0$, in the first case we have

$$\begin{aligned} a - b &= 2^{n-k} + \sum_{i=k+1}^n [(-1)^{1+\sigma_i} a_i 2^{n-i}] - \sum_{i=k+1}^n [(-1)^{1+\sigma'_i} b_i 2^{n-i}] \geq \\ &\geq 2^{n-k} - 2^{n-k-1} - 2^{n-k-1} = 0 \end{aligned}$$

because $\min \sum_{i=k+1}^n [(-1)^{1+\sigma_i} a_i 2^{n-i}] = -2^{n-k-1}$ and

$$\max \sum_{i=k+1}^n [(-1)^{1+\sigma'_i} b_i 2^{n-i}] = 2^{n-k-1}.$$

In the second case, the difference is

$$\begin{aligned} a - b &= -2^{n-k} + \sum_{i=k+1}^n [(-1)^{1+\sigma_i} a_i 2^{n-i}] - \sum_{i=k+1}^n [(-1)^{1+\sigma'_i} b_i 2^{n-i}] \leq \\ &\leq -2^{n-k} + 2^{n-k-1} + 2^{n-k-1} \end{aligned}$$

because $\min \sum_{i=k+1}^n [(-1)^{1+\sigma_i} a_i 2^{n-i}] = -2^{n-k-1}$ and

$$\max \sum_{i=k+1}^n [(-1)^{1+\sigma'_i} b_i 2^{n-i}] = 2^{n-k-1}.$$

Note that the extreme values in both cases are achieved only under the condition

$$1 - a_{k+1} = a_{k+2} = \dots = a_n = 0 = 1 - b_{k+1} = b_{k+2} = \dots = b_n. \quad \square$$

3. G -representation of the fractional part of a real number

Theorem 4 ([7]). *For any number $x \in [0; \frac{1}{2}]$, there exists a sequence of zeros and ones (α_n) such that*

$$x = \frac{\alpha_1}{2} + \sum_{k=2}^{\infty} \frac{\alpha_k (-1)^{\sigma_k}}{2^k} = \frac{\alpha_1}{2} + \sum_{k=2}^{\infty} \frac{\alpha_k}{2^{k-\sigma_k} (-2)^{\sigma_k}} \equiv \Delta_{\alpha_1 \alpha_2 \dots \alpha_n \dots}^G, \quad (2)$$

where $\sigma_k = \alpha_1 + \alpha_2 + \dots + \alpha_{k-1}$.

Corollary 1. *For any number $x \in [0; 1]$, there exists a sequence of zeros and ones (α_n) such that*

$$x = \frac{1}{2} \alpha_0 + \sum_{k=1}^{\infty} \frac{\alpha_k (-1)^{1+\sigma_k}}{2^k} \equiv \Delta_{\alpha_0 \alpha_1 \dots \alpha_n \dots}.$$

The symbolic notation $\Delta_{\alpha_1 \alpha_2 \dots \alpha_n \dots}^G$ is called the G -representation of a number $x \in [0; \frac{1}{2}]$ and of its expansion in series (2). Most numbers in the interval $[0; \frac{1}{2}]$ have a unique G -representation, and they are called G -unary numbers.

A countable set of numbers has exactly two G -representations:

$$\Delta_{c_1 \dots c_{m-1} 01(0)}^G = \Delta_{c_1 \dots c_{m-1} 11(0)}^G,$$

and such numbers are called G -binary numbers. The numbers $0 = \Delta_{(0)}^G$ and $0,5 = \Delta_{1(0)}^G$ have a unique representation and are therefore not G -binary.

Evidently, every G -binary number is rational. However not every rational number is G -binary. This is illustrated by the example:

$$\Delta_{(10)}^G = \frac{1}{2} - \frac{1}{2^3} + \frac{1}{2^5} - \frac{1}{2^7} + \dots = \frac{2}{3}.$$

Definition 3. A G -cylinder of rank m with base $c_1 c_2 \dots c_m$ is the set

$$\Delta_{c_1 c_2 \dots c_m}^G = \{x : x = \Delta_{\alpha_1 \alpha_2 \dots}^G, \alpha_i = c_i, i = \overline{1, m}\}.$$

A cylinder is a segment $[a; b]$ with endpoints

$$a = \frac{\alpha_1}{2} + \sum_{k=2}^m \frac{(-1)^{\alpha_1 + \dots + \alpha_{k-1}} \alpha_k}{2^k}, \quad b = a + \frac{1}{2^m}.$$

Cylinders of the same rank do not overlap, and the following equality holds: $\Delta_{c_1 \dots c_m}^G = \Delta_{c_1 \dots c_m 0}^G \cup \Delta_{c_1 \dots c_m 1}^G$ along with the basic metric ratio: $2|\Delta_{c_1 \dots c_m i}^G| = |\Delta_{c_1 \dots c_m}^G|$.

For any sequence (c_n) , the following holds: $\bigcap_{m=1}^{\infty} \Delta_{c_1 \dots c_m}^G = \Delta_{c_1 \dots c_m}^G$, which justifies considering a point as a cylinder of infinite rank.

It is easy to prove the following statement: if the G -representation of a number x is periodic, then x is rational.

4. Left shift and right shift operators

Remark 3. For the following functions to be well defined, we agree that among the two representations of the same G -binary number, we will use the representation $\Delta_{c_1 \dots c_{m-1} 01(0)}^G$.

Definition 4. The operator ω^n of an n -fold left shift of the digits in the G -representation of numbers is the mapping defined by the equation

$$\omega^n(x = \Delta_{\alpha_1 \alpha_2 \dots \alpha_n \dots}^G) = \Delta_{\alpha_{n+1} \alpha_{n+2} \dots \alpha_{n+k} \dots}^G.$$

The operator ω^n has the following analytic expression:

$$\omega^n(x) = 2^n (-1)^{\sigma_n} x - \alpha_1 \cdot 2^{n-1} + (-1)^{\alpha_1-1} \alpha_2 2^{n-2} + \dots + (-1)^{\alpha_1 + \dots + \alpha_{n-1}-1} 2.$$

For any $n \in N$, the operator ω^n is a piecewise-linear continuous function that attains its maximum value of 0.5 or minimum value of 0 at each G -binary point. This fundamentally distinguishes this representation from all other known two-symbol systems of encoding of numbers.

Lemma 1. *The equation $\omega^n(x) = x$ has 2^n solutions in the form $x = \Delta_{(\alpha_1 \dots \alpha_n)}^G$, where $\alpha_1, \dots, \alpha_n$ are independent variables that take values 0 and 1.*

Proof. Indeed, if $x = \Delta_{\alpha_1 \dots \alpha_n \dots}^G$ is a solution to the equation, then according to the definition of the operator $y = \omega^n(x)$ we have

$$\alpha_1 = \alpha_{n+1}, \dots, \alpha_n = \alpha_{2n}, \alpha_{n+1} = \alpha_{2n+1} = \alpha_1, \dots, \alpha_{2n} = \alpha_{3n} = \alpha_n.$$

So, $x = \Delta_{(\alpha_1 \alpha_2 \dots \alpha_n)}^G$. And the fact that $\Delta_{(\alpha_1 \dots \alpha_n)}^G$, where $\alpha_k \in A$ are free variables, is the root of the equation is obvious.

Since the image of a G -binary number under the mapping ω^n is a G -binary number of lower rank, and G -binary numbers of different ranks are never equal, we see that roots of the equation $\omega^n(x) = x$ cannot be a G -binary number. $\square$

Definition 5. The operator of the right shift of the digits in a G -representation with prefix $i_1 \dots i_k$ is the mapping $\tau_{i_1 \dots i_k}$, defined by the equation

$$\tau_{i_1 i_2 \dots i_k}(x = \Delta_{\alpha_1 \alpha_2 \dots \alpha_n \dots}^G) = \Delta_{i_1 i_2 \dots i_k \alpha_1 \alpha_2 \dots \alpha_n \dots}^G.$$

Each such function is linear and has the following analytic expression:

$$\tau_{i_1 i_2 \dots i_k}(x = \Delta_{\alpha_1 \alpha_2 \dots \alpha_n \dots}^G) = \frac{(-1)^{i_1 + \dots + i_k}}{2^k} x + \sum_{m=1}^k \frac{i_m (-1)^{i_1 + \dots + i_{m-1}}}{2^m}.$$

The following equalities are evident:

$$\omega^n(\tau_{i_1 \dots i_n}(x)) = x, \quad \tau_{\alpha_1(x) \dots \alpha_n(x)}(\omega^n(x)) = x.$$

Theorem 5. *The equation $\omega^n(x) = \tau_{i_1 \dots i_k}(x)$ has 2^n solutions in the form $x = \Delta_{(\alpha_1 \dots \alpha_n i_1 \dots i_k)}^G$, where α_j are independent variables that take values 0 and 1.*

Proof. Let $x = \Delta_{\alpha_1 \alpha_1 \dots \alpha_n \dots}^G$ be a solution of the equation. Then

$$\begin{aligned} \alpha_{n+1} &= i_1, & \dots & & \alpha_{n+k} &= i_k; \\ \alpha_{n+k+1} &= \alpha_1, & \dots & & \alpha_{n+k+n} &= \alpha_n; \\ \alpha_{2n+k+1} &= \alpha_{n+1} = i_1, & \dots & & \alpha_{2n+k+k} &= \alpha_{n+k} = i_k \text{ and so on.} \end{aligned}$$

Thus, $x = \Delta_{(\alpha_1 \dots \alpha_n i_1 \dots i_k)}^G$. It is evident that the number $\Delta_{(\alpha_1 \dots \alpha_n i_1 \dots i_k)}^G$ is a solution to the equation. $\square$

Corollary 2. *Each solution of the equation $\omega^n(x) = \tau_{i_1 \dots i_k}(x)$ is a G -unary number and, therefore, an irrational number.*

Theorem 6. *The system of equations $\omega^n(x) = x = \tau_{i_1 \dots i_k}(x)$ has a unique solution $x = \Delta_{(i_1 \dots i_k)}^G$ if $n \vdots k$.*

Proof. Let $n = k$, then the equation $\tau_{i_1 \dots i_k}(x) = x$ has a unique solution $x = \Delta_{(i_1 \dots i_k)}^G$, which obviously satisfies $\omega^k(x) = x$ as well.

According to Theorem 5, the equation $\omega^n(x) = \tau_{i_1 \dots i_k}(x)$ has solutions in the form $x = \Delta_{(\alpha_1 \dots \alpha_n i_1 \dots i_k)}^G$ while the equation $\tau_{i_1 \dots i_k}(x) = x$ has a unique solution $x = \Delta_{(i_1 \dots i_k)}^G$. Thus, when n is a multiple of k , the unique solution of the system is $x = \Delta_{(i_1 \dots i_k)}^G$. $\square$

Corollary 3. *The solution of the system of equations $\omega^k(x) = x = \tau_{i_1 \dots i_k}(x)$ is an irrational number.*

Definition 6. The G -representations of the numbers $x_1 = \Delta_{\alpha_1 \alpha_2 \dots \alpha_n \dots}^G$ and $x_2 = \Delta_{c_1 c_2 \dots c_n \dots}^G$ are said to have the same tail if there exist indices k and m such that $\alpha_{k+j} = c_{m+j}$ for all $j \in \mathbb{N}$. This is denoted symbolically as $x_1 \sim x_2$.

The binary relation “to have the same tail” is an equivalence relation. The set of all representations of a number that share the same tail, i.e., an element of the factor set, is called a tail set. Each tail set is countable, whereas the set of all tail sets has the cardinality of the continuum. All G -binary numbers form a single tail set.

A function f , defined on the interval $[0; 0, 5]$, is said to preserve the tails of the G -representations of numbers if for any $x \in [0; 0, 5]$, it holds that $f(x) \sim x$.

Recall that a bijective (i.e., one-to-one and onto) mapping of a set onto itself is called a transformation of this set. It is known [9] that the set of all continuous transformations of the interval $[0; 0, 5]$ that preserve the tails of G -representations forms an infinite, noncommutative group under the composition of transformations. Examples of such transformations include

$$f_1(x) = \begin{cases} \tau_1(x) & \text{if } 0 \leq x \leq x_1 = \Delta_{(1)}^G, \\ \omega(x) & \text{if } x_1 \leq x \leq 0, 5; \end{cases}$$

$$f_2(x) = \begin{cases} \tau_1(x) & \text{if } 0 \leq x \leq x_2 = \Delta_{(101)}^G, \\ \omega(x) & \text{if } x_2 \leq x \leq 0, 5. \end{cases}$$

5. G_2 -representation of numbers of interval $[0, g_0]$

G -representation is a special case of a more general two-symbol representation with two bases of different signs. Let us recall its definition. Suppose two bases are fixed: $g_0 \in (0; \frac{1}{2}]$, $g_1 \equiv g_0 - 1$, and the numbers $\delta_0 = 0$, $\delta_1 = g_0$.

Theorem 7 ([8]). *For any number $x \in [0; g_0]$, there exists a sequence of zeros and ones (α_n) such that*

$$x = \delta_{\alpha_1} + \sum_{k=2}^{\infty} (\delta_{\alpha_k} \prod_{j=1}^{k-1} g_{\alpha_j}) \equiv \Delta_{\alpha_1 \alpha_2 \dots \alpha_k \dots}^{G_2}, \delta_{\alpha_k} = \alpha_k g_{1-\alpha_k}. \quad (3)$$

Corollary 4. *If $g_0 = \frac{1}{2}$, then $\delta_{\alpha_k} = \frac{\alpha_k}{2^k}$, $\prod_{j=1}^{k-1} g_{\alpha_j} = \frac{(-1)^{\alpha_1 + \dots + \alpha_{k-1}}}{2^{k-1}}$ and series (3) takes the form of (2).*

The symbolic notation $\Delta_{\alpha_1 \alpha_2 \dots \alpha_k \dots}^{G_2}$ of the number x and its expansion in the alternating series (3) is called the G_2 -representation, and α_k is called its k th digit.

Lemma 2. *The Lebesgue measure of the set*

$C \equiv C[G_2; \overline{s_1 \dots s_m}] = \{x : x = \Delta_{\alpha_1 \dots \alpha_n \dots}^{G_2}, \overline{\alpha_k \dots \alpha_{k+m-1}} \neq \overline{s_1 \dots s_m}, k \in N\}$
of numbers from the interval $[0, g_0]$, whose G_2 -representation does not contain the sequence $s_1 s_2 \dots s_m$ as consecutive digits, is equal to zero.

Proof. Consider the set

$$E = \{x : x = \Delta_{\alpha_1 \alpha_2 \dots \alpha_n \dots}^{G_2}, \overline{\alpha_{km+1} \dots \alpha_{km+m}} \neq \overline{s_1 \dots s_m} \quad \forall k \in N\}.$$

Clearly, $C \subset E$. We will prove that the Lebesgue measure $\lambda(E) = 0$, which will imply that $\lambda(C) = 0$. Define $F_0 \equiv [0; g_0]$, and let F_k be the union of all cylinders of rank km that contain points from E as interior points. Let $\overline{F}_{k+1} \equiv F_k \setminus F_{k+1}$. Then, we have $F_k = F_{k+1} \cup \overline{F}_{k+1}$, $\lambda(F_{k+1}) = \lambda(F_k) - \lambda(\overline{F}_{k+1})$, $E \supset F_k \supset F_{k+1} \quad \forall k \in N$ and $\lambda(E) = \lim_{k \rightarrow \infty} \lambda(F_k)$.

Now, we express $\lambda(F_k)$ in the following form:

$$\begin{aligned} \lambda(F_k) &= \frac{g_0 \lambda(F_k)}{\lambda(F_{k-1})} \cdot \frac{\lambda(F_{k-1})}{\lambda(F_{k-2})} \cdot \dots \cdot \frac{\lambda(F_1)}{\lambda(F_0)} = \\ &= g_0 \prod_{i=1}^k \frac{\lambda(F_i)}{\lambda(F_{i-1})} = g_0 \prod_{i=1}^k \frac{\lambda(F_i) - \lambda(\overline{F}_i)}{\lambda(F_{i-1})} = g_0 \prod_{i=1}^k (1 - \frac{\lambda(\overline{F}_i)}{\lambda(F_{i-1})}). \end{aligned}$$

From this, it follows that

$$\lambda(E) = \lim_{k \rightarrow \infty} \lambda(F_k) = g_0 \prod_{k=1}^{\infty} \frac{\lambda(F_k)}{\lambda(F_{k-1})} = g_0 \prod_{k=1}^{\infty} \left(1 - \frac{\lambda(\overline{F}_k)}{\lambda(F_{k-1})}\right).$$

Since $|\Delta_{\alpha_1 \dots \alpha_k}^{G_2}| = |g_i| |\Delta_{\alpha_1 \dots \alpha_k}^{G_2}|$, we have $c_1 \leq \frac{|\Delta_{\alpha_1 \dots \alpha_k}^{G_2}|}{|\Delta_{\alpha_1 \dots \alpha_k}^{G_2}|} \leq c_2$, where $c_1 = \min\{g_0; -g_1\}$, $c_2 = \max\{g_0; -g_1\}$, and thus

$$0 < c_1^m \leq \frac{|\Delta_{\alpha_1 \dots \alpha_{km} s_1 \dots s_m}^{G_2}|}{|\Delta_{\alpha_1 \dots \alpha_{km}}^{G_2}|} = \prod_{i=1}^m |g_{s_i}| \leq c_2^m < 1.$$

Hence, the ratio $\frac{\lambda(\overline{F}_k)}{\lambda(F_{k-1})}$ is bounded away from zero, which means that the difference $1 - \frac{\lambda(\overline{F}_k)}{\lambda(F_{k-1})}$ is bounded away from one. Therefore, the necessary condition for the convergence of the infinite product is not satisfied, implying that the product diverges to zero. $\square$

Corollary 5. *The set C is a Cantor-type null set with a self-similar structure.*

Theorem 8. *Almost every number in the interval $[0; g_0]$ contains every possible sequence of digits in its G_2 -representation infinitely many times.*

Proof. Let $(s_1, \dots, s_m)$ be an arbitrary ordered sequence of zeros and ones, let H be the set of all numbers in $[0; g_0]$ whose G_2 -representation contains the sequence of digits $s_1 s_2 \dots s_m$ infinitely many times as consecutive digits, and let D be the set of all numbers that contain this sequence only a finite number of times as consecutive digits in their G_2 -representation.

We will prove that H is a full Lebesgue measure set, meaning that $\lambda(H) = g_0$. To do this, it is sufficient to show that $\lambda(\overline{H}) = 0$, where $\overline{H} = [0; g_0] \setminus H$. Clearly, $\overline{H} = \bigcup_{n=1}^{\infty} D_n$, where

$$D_n = \{x : x = \Delta_{\alpha_1 \dots \alpha_n}^{G_2}, \alpha_{k+1} \dots \alpha_{k+m} \neq s_1 \dots s_m, \forall k \geq n\}.$$

Now, we compute $\lambda(D_n)$. Since $D_n = \bigcup_{\alpha_1 \in A} \dots \bigcup_{\alpha_n \in A} [\Delta_{\alpha_1 \dots \alpha_n}^{G_2} \cap D]$,

$$\Delta_{\alpha_1 \dots \alpha_n}^{G_2} \cap D = \Delta_{\alpha_1 \dots \alpha_n}^{G_2} \cap H,$$

by the previous lemma, we have $\lambda(\Delta_{\alpha_1 \dots \alpha_n}^{G_2} \cap D) = 0$, and thus, $\lambda(D_n) = 0$. Hence $\lambda(\overline{H}) = 0$ because $\lambda(\overline{H}) \leq \sum_{n=1}^{\infty} \lambda(D_n) = 0$, which completes the proof. $\square$

6. Applications of G_2 -representation in the theory of locally complicated functions

Theorem 9. *If a G_2 -representation of numbers is defined by the parameter $g_0 \in (0; \frac{1}{2}]$, and $r_0 \in (0; \frac{1}{2}]$, then the system of functional equations*

$$\begin{cases} f(g_0x) = r_0f(x), \\ f(g_0 + (g_0 - 1)x) = r_0 + (r_0 - 1)f(x) \end{cases}$$

in the class of continuous functions defined on the interval $(0, g_0]$ has a unique solution. Moreover, when $g_0 = r_0$, the solution is $f(x) = x$, and when $g_0 \neq r_0$, the function $f(x)$ is singular, meaning that its derivative is equal to zero almost everywhere (with respect to the Lebesgue measure).

Proof. Let $x = \Delta_{\alpha_1 \dots \alpha_n \dots}^{G_2}$ be an arbitrary number in the interval $[0, g_0]$. Then $g_0x = g_0\delta_{\alpha_1} + g_0\delta_{\alpha_2}g_{\alpha_1} + g_0\delta_{\alpha_3}g_{\alpha_1}g_{\alpha_2} + \dots = \Delta_{0\alpha_1\alpha_2\dots}^{G_2}$,

$$g_0 + (g_0 - 1)x = g_0 + g_1x = g_0 + g_1\delta_{\alpha_1} + g_1\delta_{\alpha_2}g_{\alpha_1} + \dots = \Delta_{1\alpha_1\alpha_2\dots}^{G_2},$$

and thus, from the system of functional equations, we obtain

$$\begin{aligned} f(\Delta_{\alpha_1\alpha_2\dots}^{G_2}) &= \alpha_1r_{1-\alpha_1} + r_{\alpha_1}f(\Delta_{\alpha_2\alpha_3\dots}^{G_2}) = \\ &= \alpha_1r_{1-\alpha_1} + r_{\alpha_1}(\alpha_2r_{1-\alpha_2} + r_{\alpha_2}f(\Delta_{\alpha_3\dots\alpha_n\dots}^{G_2})) = \\ &= \alpha_1r_{1-\alpha_1} + r_{\alpha_1}\alpha_2r_{1-\alpha_2} + r_{\alpha_1}r_{\alpha_2}f(\Delta_{\alpha_3\dots\alpha_n\dots}^{G_2}) = \dots = \\ &= \alpha_1r_{1-\alpha_1} + \sum_{k=2}^n \alpha_kr_{1-\alpha_k} \prod_{i=1}^{k-1} r_{\alpha_i} + \left(\prod_{i=1}^k r_{\alpha_i} \right) f(\Delta_{\alpha_{n+1}\alpha_{n+2}\dots}^{G_2}). \end{aligned}$$

The product $\prod_{i=1}^n r_{\alpha_i} \rightarrow 0$ as $n \rightarrow \infty$ and f is continuous on $[0, g_0]$, so it attains a minimum and maximum value. Therefore, the remainder term in the expansion tends to zero, and we obtain

$$f(x) = \alpha_1r_{1-\alpha_1} + \sum_{k=2}^{\infty} \alpha_kr_{1-\alpha_k} \prod_{i=1}^{k-1} r_{\alpha_i} \equiv \Delta_{\alpha_1\alpha_2\dots}^{G'_2},$$

where the last symbolic notation is the G'_2 -representation with bases r_0 and $r_1 = r_0 - 1$.

As we see, f acts as a projector from the G_2 -representation to the G'_2 -representation of numbers, which implies the continuity and strict

monotonicity of f . When $g_0 = r_0$, the G_2 -representation and G'_2 -representation are identical, and thus, $f(x) = x$.

Now we prove that f is singular for $g_0 \neq r_0$. Since f is continuous and monotonic, by Lebesgue's theorem, it has a finite derivative almost everywhere (with respect of the Lebesgue measure). Let $C \subset [0; g_0]$ be the set of points where f has a finite derivative. If f has a derivative at a G_2 -unary point $x_0 \in [0; g_0]$, then it is expressed as

$$\begin{aligned} f'(x_0) &= \lim_{n \rightarrow \infty} \frac{|\Delta_{\alpha_1 \alpha_2 \dots \alpha_n}^{G'_2}|}{|\Delta_{\alpha_1 \alpha_2 \dots \alpha_n}^{G_2}|} = \lim_{n \rightarrow \infty} \frac{r_0 \prod_{i=1}^n |r_{\alpha_i}|}{g_0 \prod_{i=1}^n |g_{\alpha_i}|} = \\ &= \lim_{n \rightarrow \infty} \frac{r_0}{g_0} \prod_{i=1}^n \left| \frac{r_{\alpha_i}}{g_{\alpha_i}} \right| = \frac{r_0}{g_0} \prod_{i=1}^{\infty} \left| \frac{r_{\alpha_i}}{g_{\alpha_i}} \right|. \end{aligned}$$

If $g_0 \neq r_0$, then the infinite product diverges to zero, because the necessary condition for the convergence of an infinite product is not satisfied. Hence, $f'(x_0) = 0$. This conclusion applies to every point x such that f has a finite derivative at x and this point belongs to the set H of all numbers satisfying the conditions of the previous theorem. Therefore, f is a singular function. $\square$

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